

# **BRAID GROUPS ON GRAPHS**

Dissertation Defense  
for

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## Configuration Spaces

Let  $X$  be an arbitrary space. A *configuration* of  $n$  strands on  $X$  is an unordered  $n$ -tuple of  $n$  distinct points in  $X$ . The *configuration space* of  $n$  points on  $X$ ,  $UC^n X$ , is the set of all possible  $n$ -strand configurations.

$$UC^n X := (\Pi^n X \setminus \text{Diag}) / S_n$$

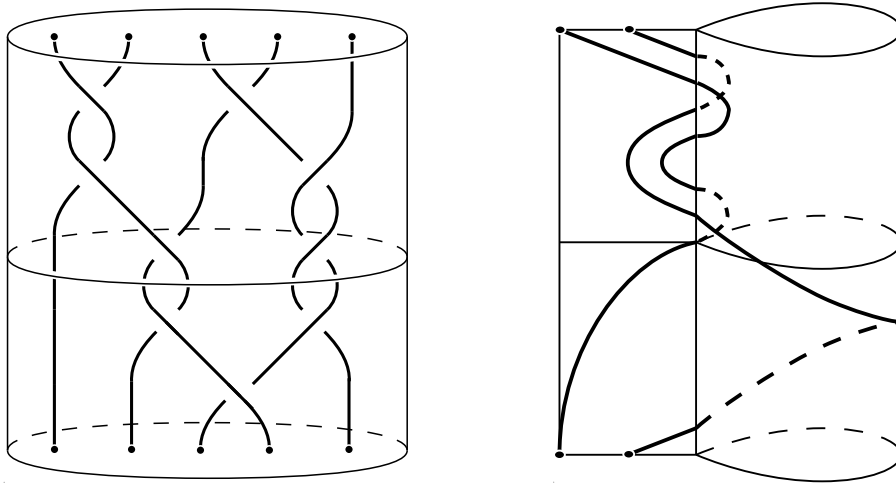
A loop in a configuration space corresponds to identical objects running around on  $X$  without colliding and returning to the initial configuration. An equivalence class of loops is a *braid*.

$$B_n X := \pi_1 UC^n X$$

## Classical and Graph Braid Groups

A classical Artin braid group is on the unit disc:

$$B_n = \pi_1 UC^n D^2$$



A graph braid group is on a graph:

$$B_n \Gamma = \pi_1 UC^n \Gamma$$

Graph braid groups can model motions of robots on factory floors.

## Discretizing Configuration Spaces

A graph  $\Gamma$  has a CW-complex structure, so  $\Pi^n \Gamma$  has a product CW-complex structure.

An open  $k$ -cell is an ordered collection of  $k$  edges and  $n - k$  vertices of  $\Gamma$ .

Define

$$Diag' = \{c \in \Pi^n \Gamma \mid c \cap Diag \neq \emptyset\}.$$

**Definition.** The *unlabelled discretized configuration space*  $UD^n \Gamma$  of  $n$  strands on  $\Gamma$  is the quotient

$$UD^n \Gamma := (\Pi^n \Gamma \setminus Diag') / S_n.$$

Under most circumstances, the unlabelled configuration space of  $\Gamma$  is homotopy equivalent to  $UD^n \Gamma$ . Specifically:

**Theorem (Abrams).** *For any  $n > 1$  and any finite graph  $\Gamma$  with at least  $n$  vertices, the unlabelled configuration space of  $n$  points on  $\Gamma$  strong deformation retracts onto  $UD^n\Gamma$  if*

- 1. each path between distinct vertices of degree not equal to 2 passes through at least  $n - 1$  edges; and*
- 2. each path from a vertex to itself that cannot be shrunk to a point in  $\Gamma$  passes through at least  $n + 1$  edges.*

Such a graph is *sufficiently subdivided* for  $n$ . We always assume sufficient subdivision.

## Some Known Results

1. (Ghrist)  $UC^n\Gamma$  is a  $K(B_n\Gamma, 1)$
2. (Ghrist)  $UC^n\Gamma \xrightarrow{sdr} X$  where  $\dim X \leq \#\{v \in \Gamma \mid \deg(v) \geq 3\}$
3. (Abrams)  $UD^n\Gamma$  is locally CAT(0), so graph braid groups have solvable word and conjugacy problems
4. (Crisp, Wiest) Graph braid groups embed nicely in right-angled Artin groups, and so are linear, biorderable, residually finite, and residually nilpotent

## Discrete Gradient Vector Fields

Let  $X$  be a finite regular CW complex, with  $K_i$  the set of open  $i$ -cells. A *discrete vector field*  $W$  on  $X$  is a sequence of partial functions  $W_i : K_i \rightarrow K_{i+1}$  such that:

1. Each  $W_i$  is injective;
2. if  $W_i(\sigma) = \tau$ , then  $\sigma < \tau$ ;
3.  $\text{im } (W_i) \cap \text{dom } (W_{i+1}) = \emptyset$ .

A *W-path* is a sequence of  $p$ -cells  $\sigma_0, \sigma_1, \dots, \sigma_r$  such that  $\sigma_{i+1} \neq \sigma_i$  and  $\sigma_{i+1} < W(\sigma_i)$ .

**Definition.** A *discrete gradient vector field* is a discrete vector field  $W$  such that  $W$  has no non-stationary closed  $W$ -paths.

A cell  $\sigma \in \text{im } W$  is *collapsible*. A cell  $\sigma \in \text{dom } W$  is *redundant*. A cell  $\sigma \in X \setminus (\text{im } W \cup \text{dom } W)$  is *critical*.

**Theorem (F, S; cf. Brown; cf. Forman).**  
*Let  $X$  be a finite regular connected CW complex with a discrete gradient vector field  $W$ . Then:*

$$\pi_1(X) \cong \langle \Sigma \mid \mathcal{R} \rangle,$$

where  $\Sigma$  is the set of positive critical 1-cells that aren't contained in  $T$ , and

$$\mathcal{R} = \{r(w) \mid w \text{ a boundary of a crit. 2-cell}\}.$$

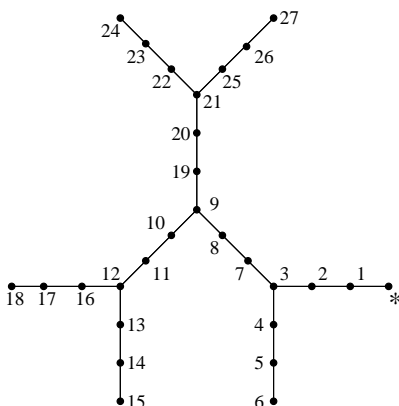
If there is a unique critical 0-cell (and there always will be), this presentation is unique given  $W$  and is called the *Morse presentation*.



## Graphs

Let  $\Gamma$  be a finite simplicial graph. If  $v \in V\Gamma$ , the degree of  $v$  is  $d(v)$ . If  $d(v) \geq 3$ ,  $v$  is *essential*. Let  $T$  be a maximal tree.

Pick  $* \in V\Gamma$  with  $d(*) = 1$ . Embed  $T$  in  $\mathbb{R}^2$ . Define  $\leq$  on  $VT = V\Gamma$  by numbering  $VT$  with a clockwise traversal of  $T$  from  $*$ .



Number edges adjacent to  $v$  from 0 in order of traversal. Each edge numbering is a *direction* from  $v$ . For  $v \neq *$ , the edge 0 from  $v$  is  $e(v)$ .

If  $e \in E\Gamma$ , the endpoints of  $e$  are  $\iota(e)$  and  $\tau(e)$ , where  $\iota(e) > \tau(e)$ . If  $e \in \Gamma \setminus T$ ,  $e$  is a *deleted edge*.

## Defining $W$ for a Graph

Let  $c \in UD^n\Gamma$ . If  $v \in c$ ,  $v$  is *blocked* in  $c$  if  $v = *$  or  $e(v) \cap c_i \neq \emptyset$  for some  $i$ . Otherwise,  $v$  is unblocked and  $(c - \{v\} + \{e(v)\})$  is a *reduction*.

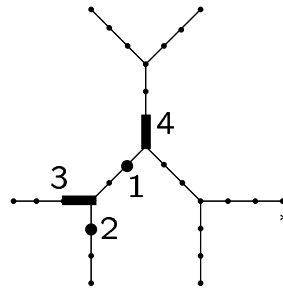
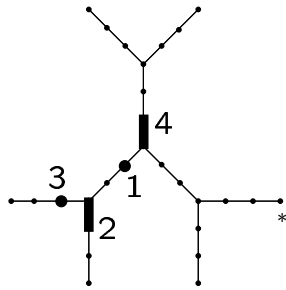
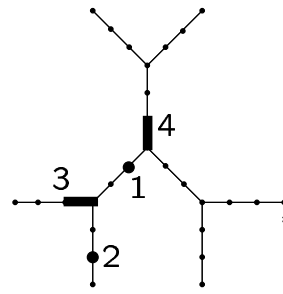
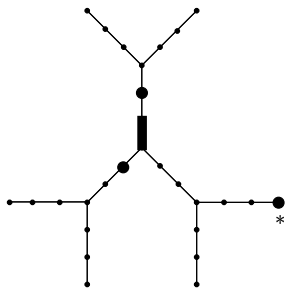
If  $e \in c$ ,  $e$  is *not order-respecting* if  $e$  is a deleted edge or  $v$  is adjacent to  $\tau(e)$  and  $\tau(e) < v < \iota(e)$  for some  $v \in c$ .

**Definition (F,S).** For  $c$  a cell of  $UD^n\Gamma$ ,

1. Traversing  $T$  from  $*$ , if an unblocked vertex is hit before an order-respecting edge,  $W(c) = (c - \{v\} + \{e(v)\})$ ;  $c$  is redundant.
2. If an order-respecting edge is first,  $c \in \text{im } W$ ;  $c$  is collapsible.
3. Otherwise,  $W(c)$  is undefined;  $c$  is critical.

**Theorem (F,S).**  $W$  is a discrete gradient vector field on  $UD^n\Gamma$ .

## Examples: The tree $T_{min}$



## Observations

**Theorem (F, S).** *Let  $\Gamma, T, W,$  be as before. Let  $D$  be the number of deleted edges. With minor assumptions, the Morse presentation for  $B_n\Gamma$  has*

$$D + \sum_{\substack{v \in V(T) \\ \text{essential}}} \sum_{i=2}^{d(v)-1} \left[ \binom{n + d(v) - 2}{n - 1} - \binom{n + d(v) - i - 1}{n - 1} \right]$$

*generators.*

**Corollary (Ghrist).** *If  $\Gamma$  is a radial tree - i.e. has exactly one essential vertex,  $v$  - then  $B_n\Gamma$  is free of rank*

$$\sum_{i=2}^{d(v)-1} \left[ \binom{n + d(v) - 2}{n - 1} - \binom{n + d(v) - i - 1}{n - 1} \right].$$

**Theorem (F, S).** *Let  $\Gamma$  be a tree and  $c$  a critical cell of  $UD^n\Gamma$ . Let*

$$k := \min \left\{ \left\lfloor \frac{n}{2} \right\rfloor, \#\{v \in \Gamma^0 \mid v \text{ is essential}\} \right\}.$$

*Then  $\dim c \leq k$ . In particular,  $UD^n\Gamma \xrightarrow{sdr} X$  where  $\dim X \leq k$ .*

**Theorem (F, S; cf. Mautner).** *Let*

$$k := \min \left\{ \left\lfloor \frac{n + 1 - \chi(\Gamma)}{2} \right\rfloor, \#\{v \in \Gamma^0 \mid v \text{ is essential}\} \right\}.$$

*$UD^n\Gamma \xrightarrow{sdr} X$  where  $\dim X \leq k$ .*

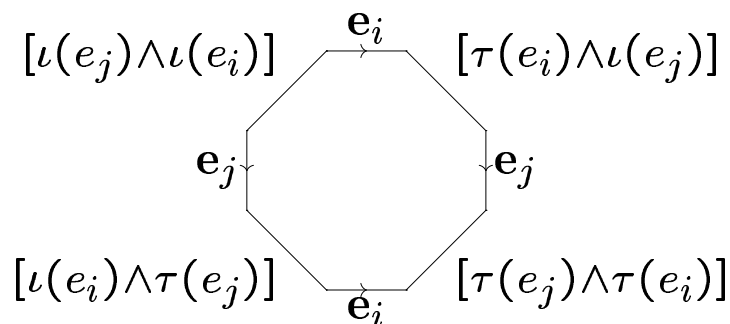
## Presentations

**Theorem (F,S).** For  $\Gamma$  a tree,  $B_n\Gamma$  has a Morse presentation where generators are critical 1-cells and relators are reduced forms  $w(c)$  of boundaries of critical 2-cells  $c$ , which are explicitly given....

A corner is an essential vertex  $A$  and  $0 < i < j < d(A)$ .

Corners  $\leftrightarrow$  critical 1-cells at essential vertices.

**Theorem (F,S).** For any  $\Gamma$ ,  $B_2\Gamma$  has a Morse presentation  $P$ , where the generators of  $P$  consist of all corners plus all deleted edges, and the relations of  $P$  are

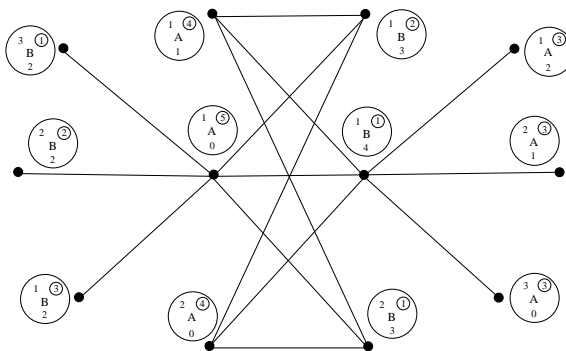


## Examples

- $B_4T_{min}$  has 24 generators and relators:

$$\begin{aligned}
 & \left[ \begin{array}{c} 1 \textcircled{3} \\ A \\ 0 \end{array}, \begin{array}{c} 1 \textcircled{1} \\ B \\ 2 \end{array} \right], \left[ \begin{array}{c} 1 \textcircled{3} \\ A \\ 0 \end{array}, \begin{array}{c} 1 \textcircled{1} \\ C \\ 2 \end{array} \right], \left[ \begin{array}{c} 1 \textcircled{3} \\ A \\ 0 \end{array}, \begin{array}{c} 1 \textcircled{1} \\ D \\ 2 \end{array} \right], \\
 & \left[ \begin{array}{c} 1 \textcircled{3} \\ B \\ 0 \end{array}, \begin{array}{c} 1 \textcircled{1} \\ D \\ 2 \end{array} \right], \left[ \begin{array}{c} 2 \textcircled{1}^{-1} \\ B \\ 1 \end{array}, \begin{array}{c} 3 \textcircled{1} \\ B \\ 0 \end{array}, \begin{array}{c} 1 \textcircled{1} \\ C \\ 2 \end{array} \right], \\
 & \left[ \begin{array}{c} 1 \textcircled{1} \\ D \\ 2 \end{array}, \begin{array}{c} 2 \textcircled{2} \\ B \\ 0 \end{array}, \begin{array}{c} 2 \textcircled{1} \\ B \\ 1 \end{array}, \begin{array}{c} 1 \textcircled{1} \\ C \\ 2 \end{array}, \begin{array}{c} 2 \textcircled{1}^{-1} \\ B \\ 1 \end{array}, \begin{array}{c} 2 \textcircled{2}^{-1} \\ B \\ 0 \end{array} \right]
 \end{aligned}$$

- $B_6T_H$ , where  $T_H$  is the H graph, is  $F_{18}$  free product the RAAG on:



- For  $\Gamma$  the 4-spoked wheel,  $B_2\Gamma \cong F_5$  (after rewriting to eliminate existing 6 relators).

## Right-Angled Artin Groups

**Definition.** Let  $\Delta$  be a finite simple graph. The *right-angled Artin group* associated to  $\Delta$  is  $G(\Delta)$  generated by  $V\Delta$ , with relations commutators corresponding to edges in  $\Delta$ .

**Theorem (Droms).** *Two graphs  $\Delta$  and  $\Delta'$  are graph isomorphic if and only if the groups  $G(\Delta)$  and  $G(\Delta')$  are group isomorphic.*

**Theorem (Crisp, Wiest).** *For any finite graph  $\Gamma$  and any  $n$ , there exists a graph  $\Delta$  such that  $B_n\Gamma$  embeds in  $G(\Delta)$ .*

**Theorem (S).** *For every finite graph  $\Delta$  and any coloring  $C$  of  $\Delta$  with  $n$  colors, there exists a graph  $\Gamma$  such that the right-angled Artin group  $G(\Delta)$  embeds into the graph braid group  $B_n\Gamma$ .*



## $\Delta$ -halos

**Definition (Halo).** A connected graph  $\Gamma$  is a  $\Delta$ -halo if for each  $a_i \in V\Delta$ , there is a simple edge loop  $\gamma_i$  in  $\Gamma$  (an *Artin loop*), such that:

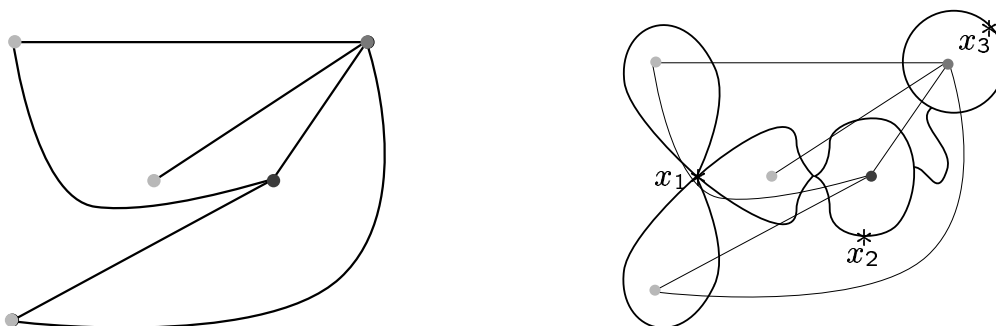
- for each color  $c \in \{1, \dots, n\}$ , there exists  $x_c \in V\Gamma$  with

$$x_c = \bigcap_{iC(a_i)=c} \gamma_i \cap \bigcap_{iC(a_i) \neq c} (\Gamma \setminus \gamma_i).$$

- if  $(a_i, a_j) \notin E(\Delta)$ , then  $\gamma_i \cap \gamma_j = v$  for some vertex  $v$ . If  $C(a_i) = C(a_j)$  then  $v = x_{C(a_i)}$ ; otherwise,  $v$  is on no other Artin loop.
- if  $(a_i, a_j) \in E(\Delta)$ ,  $\gamma_i \cap \gamma_j = \emptyset$ .

$\{x_1, \dots, x_n\} \in UC^n\Gamma$  is the *Artin basepoint*.

## Example and Sketch of Proof



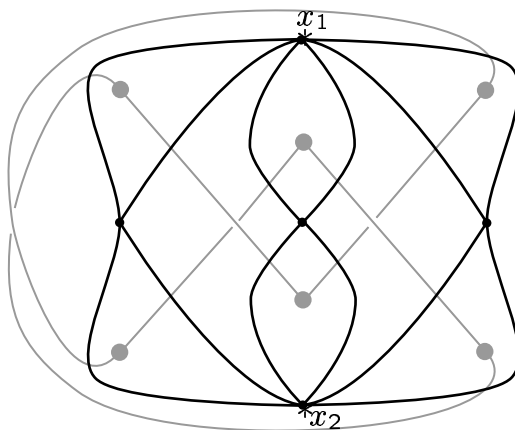
Sketch of proof of theorem: Take the element corresponding to  $a_i \in V\Delta$  to  $\gamma_i^2$ .

Compose the induced map with Crisp and Wiest's embedding back into right-angled Artin groups. Show the composition is injective using HNN structure on right-angled Artin groups and Britton's Lemma.

## Corollaries

**Corollary (S).** *Let  $n_0$  be the chromatic number for  $\Delta$ . Then there exists an embedding of  $G(\Delta)$  into  $B_{n_0}\Gamma$  for some  $\Delta$ -halo  $\Gamma$ .*

**Example.** Let  $A = G(C_6)$ . A  $C_6$ -halo for the 2-coloring is:



**Theorem (S).** *There exists a planar graph braid group which contains a hyperbolic surface subgroup, with only  $n = 2$  strands.*

## Definitions for Cohomology on Trees

Define an equivalence relation  $\sim$  on cells of  $UD^nT$ :  $c \sim c'$  if

1.  $E(c) = E(c')$ , and

2. for any connected component  $C$  of  $T - \bigcup_{e \in E(c)} e$ ,

$$|C \cap (c \setminus E(c))| = |C \cap (c' \setminus E(c'))|.$$

Define a partial order  $\leq$  on classes:  $[c] \leq [c_1]$  if there exists  $\hat{c} \in [c]$ ,  $\hat{c}_1 \in [c_1]$  such that  $\hat{c} \leq \hat{c}_1$ .

Define the cellular cocycle  $\phi_{[c]}$  to be:

$$\begin{aligned} \phi_{[c]}(\tilde{c}) &= 1 && \text{if } \tilde{c} \sim c \\ \phi_{[c]}(\tilde{c}) &= 0 && \text{otherwise.} \end{aligned}$$

## Homology and Cohomology on Trees

**Theorem (F).**  $H_i(UD^nT; \mathbb{Z}/2\mathbb{Z})$  is free abelian with basis  $\{c \mid c \text{ a critical } i\text{-cell}\}$ .

**Theorem (F,S).**

1.  $H^i(UD^nT; \mathbb{Z}/2\mathbb{Z})$  is free abelian with basis  $\{c^* \mid c \text{ a critical } i\text{-cell}\}$ .

2. For  $c$  critical, if  $[c]$  is the l.u.b. of distinct classes  $[c_1], \dots, [c_i]$  of 1-cells, then, WLOG,  $c_1, \dots, c_i$  are critical and

$$c_1^* \cup \dots \cup c_i^* = c^*.$$

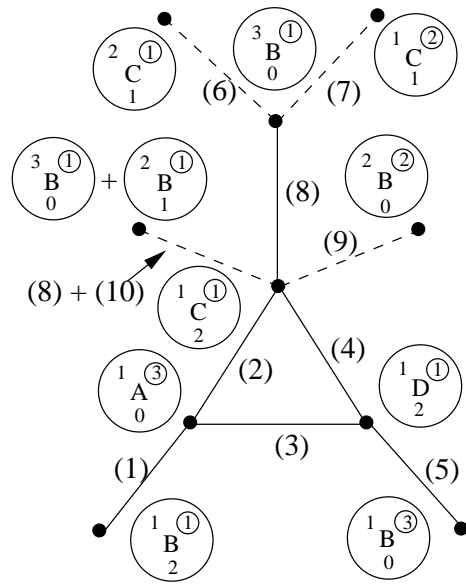
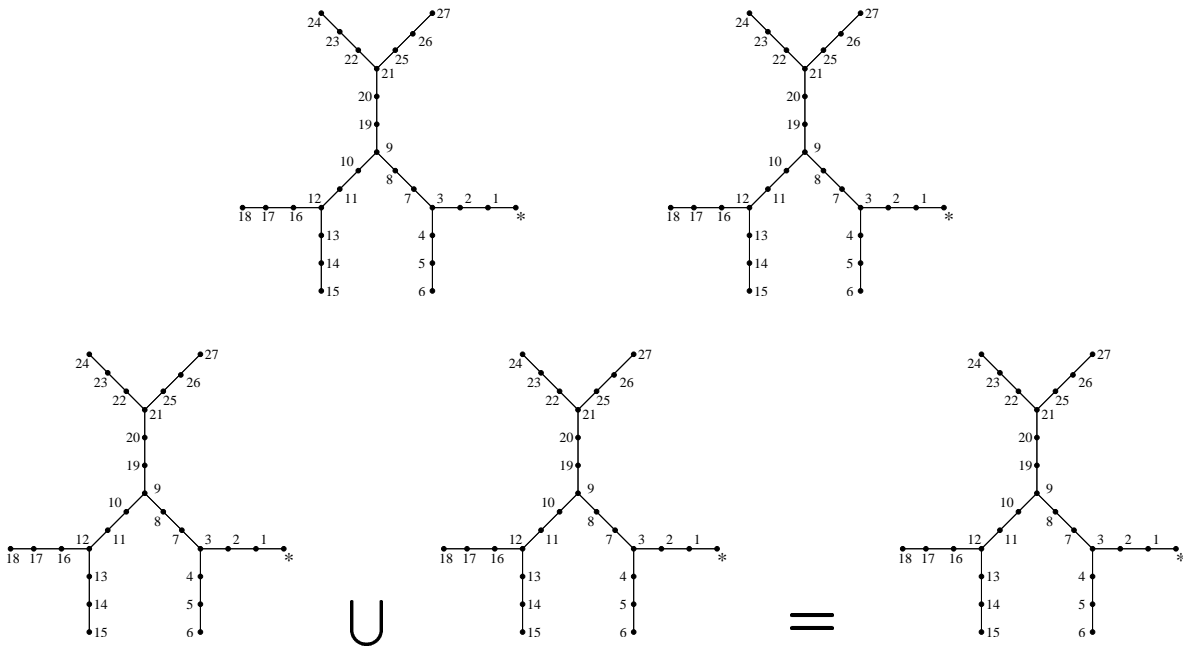
3. If  $[c_1], \dots, [c_i]$  are distinct classes of 1-cells having l.u.b.  $[c]$ , then

$$[\phi_{[c_1]}] \cup \dots \cup [\phi_{[c_i]}] = [\phi_{[c]}].$$

Otherwise,

$$[\phi_{[c_1]}] \cup \dots \cup [\phi_{[c_i]}] = 0.$$

# Example: $T_{min}$



## Ghrist's Conjecture

**Conjecture (Ghrist; Abrams).** **Ghrist's Conjecture.** *The (pure) braid group of any planar graph is a right-angled Artin group.*

**Theorem (Abrams).**  *$B_2K_5$  and  $B_2K_{3,3}$  are hyperbolic surface groups, and are therefore not right-angled Artin.*

**Theorem (Connolly, Doig).** *Braid groups on linear trees are right-angled Artin.*

- There are known examples of tree braid groups which are not right-angled Artin (Mautner).

## Exterior Face Algebras

**Definition.** Let  $K$  be a finite simplicial complex with vertices  $\{v_1, \dots, v_n\}$ . The  $(\mathbb{Z}/2\mathbb{Z})$  exterior face algebra  $\Lambda(K)$  on  $K$  has:

- basis  $\{v_{i_1} \dots v_{i_j} \mid i_1 < \dots < i_j\}$

- relations

$$v_i v_j = v_j v_i$$

$$v_i^2 = 0$$

$$v_{i_1} \dots v_{i_k} = 0 \quad \text{when } i_1 < \dots < i_k \text{ and } \{v_{i_1}, \dots, v_{i_k}\} \notin K.$$

**Theorem (Gubeladze).**  $\Lambda(K) \cong \Lambda(K')$  if and only if  $K \cong K'$ .



## Counterexamples to Ghrist's Conjecture

**Theorem (F,S).** *Let  $T$  be a finite tree. The tree braid group  $B_n T$  is a right-angled Artin group if and only if  $T$  is linear or  $n < 4$ .*

Proof relies on:

- Cohomology for any RAAG is an exterior face algebra over a flag complex.
- $H^*(T_{min})$  is an exterior face algebra, but is not over a flag complex.
- Embedding  $T_{min}$  into a nonlinear tree induces nice maps on cohomology.

## Cohomology for $B_4T$

**Theorem (S).** *Let  $T$  be a tree. For  $n = 4$ ,  $H^*(B_4T; \mathbb{Z}/2\mathbb{Z})$  is an exterior face algebra. The simplicial complex  $\Delta$  defining the exterior face algebra structure is unique and 1-dimensional.*

Idea of proof: usually, if  $c_1$  and  $c_2$  are critical, then

$$c_1^* \cup c_2^* = c^*$$

where  $c$  is critical.

Take care of the 'usually': rewrite Morse basis so exceptional cases no longer cup with anything, and leave everything else alone.

## Rigidity for $B_4T$

**Theorem (S).** *Two finite trees are homeomorphic if and only if their 4 strand braid groups are isomorphic.*

Idea of proof: Reconstruct  $T$  from  $\Delta$  using neighborhood hierarchies.

$$\begin{array}{ccc} T & \xrightarrow{\quad} & B_n T \\ & \searrow \psi_n & \downarrow \\ \Delta & \xleftarrow{\quad} & H^* B_n T \end{array}$$

## Neighborhood Hierarchies

Let  $N_v := \{v' \in \Delta \mid (v, v') \in EV\}$  be the *neighborhood* of  $v$ . Subset inclusion defines a pre-order:  $v \leq_N v' \leftrightarrow \emptyset \subsetneq N_v \subseteq N_{v'}$ .

For any  $[v_0]$ , form a  $\leq_N$ -hierarchy  $H_{v_0}$ : a graph with vertices  $[v]$  for  $[v] \leq_N [v_0]$  and edges for largest subsets.

The graph  $H_{v_0}$  is a tree, a subtree of  $T$ , and just needs more leaves.

## The Isomorphism Problem for $B_4T$

**Theorem (S).** *There exists an algorithm which, given two four strand tree braid groups on finite trees, decides whether the two groups are isomorphic.*

Also true for  $n = 5$  or  $n \geq 4$ ?